# Interpolation in exchange-free logics 

Wesley Fussner<br>Mathematical Institute, University of Bern<br>Switzerland<br>Joint work with G. Metcalfe and S. Santschi

CLoCk 68
Cracow, Poland

28 June 2023

## Interpolation generally

- This talk is about interpolation, which deals with certain kinds of explanations for why given inferences hold.
- Craig interpolation property (CIP): If $\vdash \varphi \rightarrow \psi$, then there exists a formula $\delta$ such that $\operatorname{var}(\delta) \subseteq \operatorname{var}(\varphi) \cap \operatorname{var}(\psi)$ and $\vdash \varphi \rightarrow \delta$ and $\vdash \delta \rightarrow \psi$.
- Various versions designed for particular applications: Uniform interpolation (databases), feasible interpolation (complexity theory), McMillan-style Craig interpolation (hardware and software verification), and so on.
- Deductive interpolation property (DIP): If $\Gamma \vdash \varphi$, then there exists a set of formulas $\Gamma^{\prime}$ such that $\operatorname{var}\left(\Gamma^{\prime}\right) \subseteq \operatorname{var}(\Gamma) \cap \operatorname{var}(\varphi)$ and $\Gamma \vdash \Gamma^{\prime}$ and $\Gamma^{\prime} \vdash \varphi$.

Broadly, interpolation is understood as a rather uncommon property.

- Exactly 7 consistent superintuitionistic logics with CIP/DIP, just 3 positive logics (Maksimova 1977).
- $\leq 38$ normal extensions of S4 with CIP.
- Uncountably many extensions of Hájek's basic fuzzy logic without DIP (Montagna 2006).
- Positive results tend to use specialized methods and be fairly limited in scope.
- Intuitionistic logic is a substructural logic.
- Generally these arise from dropping/relaxing some of the structural rules appearing in Gentzen's sequent calculus presentation of intuitionistic logic (exchange, weakening, contraction).
- Substructural logics encompass many logics arising independently:
- Hájek's basic fuzzy logic and Łukasiewicz logic
- The most prominent relevant logics
- Linear logic and bunched implication logics
- Substructural logics can be formulated under the umbrella of extensions of the full Lambek calculus.

Identity Axioms
$\overline{\alpha \Rightarrow \alpha}{ }^{(\text {ID })}$
Left Operation Rules
$\frac{\Gamma_{1}, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \mathrm{e}, \Gamma_{2} \Rightarrow \Delta}(\mathrm{e} \Rightarrow)$
$\overline{\mathrm{f} \Rightarrow}(\mathrm{f} \Rightarrow)$
$\frac{\Gamma_{2} \Rightarrow \alpha \quad \Gamma_{1}, \beta, \Gamma_{3} \Rightarrow \Delta}{\Gamma_{1}, \beta / \alpha, \Gamma_{2}, \Gamma_{3} \Rightarrow \Delta}(/ \Rightarrow)$
$\frac{\Gamma_{2} \Rightarrow \alpha \quad \Gamma_{1}, \beta, \Gamma_{3} \Rightarrow \Delta}{\Gamma_{1}, \Gamma_{2}, \alpha \backslash \beta, \Gamma_{3} \Rightarrow \Delta}(\backslash \Rightarrow)$
$\frac{\Gamma_{1}, \alpha, \beta, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \alpha \cdot \beta, \Gamma_{2} \Rightarrow \Delta}(\cdot \Rightarrow)$
$\frac{\Gamma_{1}, \alpha, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \alpha \wedge \beta, \Gamma_{2} \Rightarrow \Delta}(\wedge \Rightarrow)_{1}$
$\frac{\Gamma_{1}, \beta, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \alpha \wedge \beta, \Gamma_{2} \Rightarrow \Delta}(\wedge \Rightarrow)_{2}$
$\frac{\Gamma_{1}, \alpha, \Gamma_{2} \Rightarrow \Delta \quad \Gamma_{1}, \beta, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \alpha \vee \beta, \Gamma_{2} \Rightarrow \Delta}(\vee \Rightarrow)$

Cut Rule

$$
\frac{\Gamma_{2} \Rightarrow \alpha \quad \Gamma_{1}, \alpha, \Gamma_{3} \Rightarrow \Delta}{\Gamma_{1}, \Gamma_{2}, \Gamma_{3} \Rightarrow \Delta}(\mathrm{cUT})
$$

Right Operation Rules

$$
\begin{aligned}
& \overline{\Rightarrow \mathrm{e}}(\Rightarrow \mathrm{e}) \\
& \frac{\Gamma \Rightarrow}{\Gamma \Rightarrow \mathrm{f}}(\Rightarrow \mathrm{f}) \\
& \frac{\Gamma, \alpha \Rightarrow \beta}{\Gamma \Rightarrow \beta / \alpha}(\Rightarrow /) \\
& \frac{\alpha, \Gamma \Rightarrow \beta}{\Gamma \Rightarrow \alpha \backslash \beta}(\Rightarrow \mathrm{V}) \\
& \frac{\Gamma_{1} \Rightarrow \alpha \quad \Gamma_{2} \Rightarrow \beta}{\Gamma_{1}, \Gamma_{2} \Rightarrow \alpha \cdot \beta}(\Rightarrow \cdot) \\
& \frac{\Gamma \Rightarrow \alpha}{\Gamma \Rightarrow \alpha \vee \beta}(\Rightarrow \vee)_{1} \\
& \frac{\Gamma \Rightarrow \beta}{\Gamma \Rightarrow \alpha \vee \beta}(\Rightarrow \vee)_{2} \\
& \frac{\Gamma \Rightarrow \alpha{ }^{2} \Rightarrow \beta}{\Gamma \Rightarrow \alpha \wedge \beta}(\Rightarrow \wedge)
\end{aligned}
$$

## Basic structural rules

$$
\begin{gathered}
\frac{\Gamma_{1}, \Pi_{1}, \Pi_{2}, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \Pi_{2}, \Pi_{1}, \Gamma_{2} \Rightarrow \Delta}(\mathrm{EL}) \\
\frac{\Gamma_{1}, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \Pi, \Gamma_{2} \Rightarrow \Delta}(\mathrm{wL}) \quad \frac{\Pi \Rightarrow}{\Gamma_{1}, \Pi, \Gamma_{2} \Rightarrow \Delta}(\mathrm{wR}) \\
\frac{\Gamma_{1}, \Pi, \Pi, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \Pi, \Gamma_{2} \Rightarrow \Delta}(\mathrm{cL}) \\
\frac{\Gamma_{1}, \Pi_{1}, \Gamma_{2} \Rightarrow \Delta \quad \Gamma_{1}, \Pi_{2}, \Gamma_{2} \Rightarrow \Delta}{\Gamma_{1}, \Pi_{1}, \Pi_{2}, \Gamma_{2} \Rightarrow \Delta}(\text { mingle })
\end{gathered}
$$

## Interpolation and exchange

Lots of success with DIP in the context of exchange without much systematic information:

- Lots of work from proof theory (Maehara, Ono, others).
 2000).
- Continuum-many extensions of FL + exchange with DIP, also for full linear logic (F.-Santschi 2023). Depends heavily on group theory.
- Previously thought that there may be no extension of FL lacking exchange with DIP (Gil Férez-Ledda-Tsinakis 2015).
- Example given in 2020 by Gil-Férez, Jipsen, Metcalfe.
- Several natural examples involving the law of excluded middle (F.-Galatos 2022).


## Today's contribution

We will see that:

- There are continuum-many axiomatic extensions of FL without exchange that have DIP.
- All have the contraction and mingle rules, and are characteristic with respect to linearly ordered models (semilinear).
- Among axiomatic extensions of falsum-free FL + contraction + mingle + exchange + semilinearity, only 60 with DIP.

Part I:
The set-up

## Logics without exchange

- Note that exchange is derivable in the presence of contraction + left weakening.
- So, if we want to study extensions of FL without exchange while doing minimal mutilation to the intuitionistic framework, we can't keep both contraction and weakening.
- Natural solution: Replace one of contraction or weakening by a slightly less powerful rule.
- Here we replace weakening by the mingle rule.
- We thus focus on $\mathrm{FL}_{c m}$, full Lambek calculus + contraction + mingle.
- We also consider the variant without falsum $f$.


## Algebraic semantics

- Key methodology: Algebraization of the consequence relation of FL.
- Algebraization gives mutually inverse, back-and-forth translations between a consequence relation and the equational consequence relation of some class of algebraic models (in our case, residuated lattices).
- Transfer many properties by bridge theorems:
- Local deduction theorems correspond to the congruence extension property.
- With the above, DIP corresponds to the amalgamation property.


## Residuated lattices

A residuated lattice is an algebraic structure of the form $(A, \wedge, \vee, \cdot, \backslash, /, e)$ where

- $(A, \wedge, \vee)$ is a lattice,
- $(A, \cdot, e)$ is a monoid, and
- for all $x, y, z \in A$,

$$
x \cdot y \leq z \Longleftrightarrow y \leq x \backslash z \Longleftrightarrow x \leq z / y
$$

We use all the expected terminology: Commutative, idempotent, totally ordered, linear, etc.

Semilinear: Subalgebra of a direct product of totally ordered residuated lattices.

Note that despite the adjunction condition, residuated lattices form a variety (equational class). Subvarieties of residuated lattices correspond exactly with axiomatic extensions of FL without falsum.

## Some corresponding properties

Because of algebraization, there's a back-and-forth dictionary of concepts:

- Exchange corresponds to commutativity $x y=y x$.
- (Left) weakening correspond to integrality $x \leq e$.
- Contraction corresponds to the square-increasing law $x \leq x^{2}$.
- Mingle corresponds to the square-decreasing law $x^{2} \leq x$.
- So, contraction + mingle corresponds to multiplication being idempotent $x^{2}=x$.
- To study axiomatic extensions of positive FL + contraction + mingle, we can study varieties (equational classes) of idempotent residuated lattices.
- Semilinearity corresponds to the communication rule.


## Amalgamation

## Definition:

Let $\mathcal{K}$ be a class of algebraic structures. A span in $\mathcal{K}$ is a quintuple $(A, B, C, f, g)$, where $A, B, C \in \mathcal{K}$ and $f: A \rightarrow B, g: A \rightarrow C$ are embeddings. We say that $\mathcal{K}$ has the amalgamation property (or $\mathrm{AP})$ if for every span $(A, B, C, f, g)$ in $\mathcal{K}$ there exists $D \in \mathcal{K}$ and embeddings $f^{\prime}: B \rightarrow D$ and $g^{\prime}: C \rightarrow D$ such that $f^{\prime} \circ f=g^{\prime} \circ g$.


## Amalgamation

## Definition:

Let $\mathcal{K}$ be a class of algebraic structures. A span in $\mathcal{K}$ is a quintuple $(A, B, C, f, g)$, where $A, B, C \in \mathcal{K}$ and $f: A \rightarrow B, g: A \rightarrow C$ are embeddings. We say that $\mathcal{K}$ has the amalgamation property (or $\mathrm{AP})$ if for every span $(A, B, C, f, g)$ in $\mathcal{K}$ there exists $D \in \mathcal{K}$ and embeddings $f^{\prime}: B \rightarrow D$ and $g^{\prime}: C \rightarrow D$ such that $f^{\prime} \circ f=g^{\prime} \circ g$.


Part II:
The case without exchange

- To get continuum-many axiomatic extensions of FL + contraction + mingle with the DIP, it's enough to come up with continuum-many varieties of semilinear idempotent residuated lattices with the amalgamation property.
- We're inspired by Galatos 2005, which gives continuum-many atoms in the lattice of subvarieties of semilinear idempotent residuated lattices (logics with no non-trivial extensions).
- We'll show that each of Galatos's varieties have the amalgamation property.
- This involves four ingredients: The nested sum construction of residuated lattices, the symbolic dynamics of bi-infinite words, tools from first-order model theory, and new characterizations of the AP.


## Starting out

Suppose $S \subseteq \mathbb{Z}$. We define an algebra on

$$
A_{S}=\left\{a_{i}: i \in \mathbb{Z}\right\} \cup\left\{b_{j}: j \in \mathbb{Z}\right\} \cup\{e\}
$$

Order the elements of $A_{S}$ by setting $b_{i}<b_{j}<e<a_{k}<a_{l}$ if and only if $i, j, k, I \in \mathbb{Z}$ with $i<j$ and $l<k$. Further, for $i, j \in \mathbb{Z}$ define $a_{i} a_{j}=a_{\min \{i, j\}}, b_{i} b_{j}=b_{\min \{i, j\}}$, and

$$
\begin{aligned}
& a_{i} b_{j}= \begin{cases}a_{i} & \text { if } i<j \text { or } i=j \in S \\
b_{j} & \text { if } i>j \text { or } i=j \notin S\end{cases} \\
& b_{j} a_{i}= \begin{cases}b_{j} & \text { if } j<i \text { or } i=j \in S \\
a_{i} & \text { if } j>i \text { or } i=j \notin S\end{cases}
\end{aligned}
$$

We stipulate that $e$ is a multiplicative identity and define residuals $\backslash$ and / in the usual way. The residuated lattice obtained in this way is denoted by $\boldsymbol{A}_{S}$ and the variety it generates is $V_{S}$.

## Bi-infinite words

## Definition:

A word over $\{0,1\}$ is a function $w: A \rightarrow\{0,1\}$, where $A$ is some subinterval of $\mathbb{Z}$. A word is finite if $|A|$ is finite and bi-infinite if $A=\mathbb{Z}$. We say that a finite word $v: A \rightarrow\{0,1\}$ is a subword of a word $w$ if there exists an integer $k$ such that $v(i)=w(i+k)$ for all $i \in A$. The characteristic function $w_{S}$ of a subset $S \subseteq \mathbb{Z}$ is an example of a bi-infinite word.

## Definition:

We define a pre-order $\sqsubseteq$ on the set of all bi-infinite words by setting $w_{1} \sqsubseteq w_{2}$ if and only if every finite subword of $w_{1}$ is a subword of $w_{2}$. For bi-infinite words $w_{1}, w_{2}$, we write $w_{1} \cong w_{2}$ if and only if $w_{1} \sqsubseteq w_{2}$ and $w_{2} \sqsubseteq w_{1}$.

## Fact:

There are continuum-many pairwise incomparable minimal bi-infinite words.

## Constructing the subvarieties

- For each $S \subseteq \mathbb{Z}$, we can consider $S$ as a bi-infinite word by identifying it with its characteristic function $w_{S}$.
- If $w_{S}$ is minimal, then $V_{S}$ gives an atom in the lattice of subvarieties of semilinear idempotent residuated lattices.
- The cardinality result for atoms follows from the fact that there continuum-many pairwise incomparable minimal bi-infinite words.
- The nested sum extends the well-known ordinal sum construction used for Hájek's basic logic.
- It is technical to state correctly, but it amounts to replacing the identity element $e$ in a residuated lattice $\mathbf{A}$ by another residuated lattice B.
- This can only be done for some residuated lattices, but it turns out that the algebra $\mathbf{A}_{S}$ are admissible.


## The key lemma

## Lemma:

Suppose that $S \subseteq \mathbb{Z}$.
(1) $\mathbb{H S P}_{U}\left(\mathbf{A}_{S}\right)$ is the class of totally ordered members of $\mathrm{V}_{S}$. In particular, $\mathbb{H S P}_{U}\left(\mathbf{A}_{S}\right)$ consists of the finitely subdirectly irreducible members of $\mathrm{V}_{S}$.
(2) If $w_{S}$ is minimal, then $\operatorname{HSP}_{U}\left(\mathbf{A}_{S}\right)$ is closed under nested sums. In particular, the finitely subdirectly irreducible members of $\mathrm{V}_{S}$ are exactly nested sums of members of $\mathrm{K}_{S}=\mathbb{I}\left(\left\{\mathbf{A}_{T}: w_{T} \sqsubseteq w_{S}\right\}\right)$.

The proof is a technical argument using ultraproducts, and invokes the fact that every algebra embeds into an ultraproduct of its finitely generated subalgebras.

## Amalgamation of chains

## Lemma:

Suppose that $S \subseteq \mathbb{Z}$ is such that $w_{S}$ is minimal. Then the class of totally ordered members in $\mathrm{V}_{S}$ has the amalgamation property.

The proof involves decomposing each chain in a given span into a nested sum of its 1-generated subalgebras (by F.-Galatos 2022), and then collecting 1 -generated subalgebras. Because the totally ordered members are closed under nested sums by the previous lemma, these can be collected into an amalgam by taking the nested sum.

## From chains upward

This doesn't quite prove that the varieties $\mathrm{V}_{S}, w_{S}$ minimal, have the AP. For this, we need to extend the AP from chains:

## Theorem (F.-Metcalfe 2022)

Suppose V is a congruence-distributive variety with the congruence extension property, and that the class of finitely subdirectly irreducibles in V is closed under taking subalgebras. Then if the class of finitely subdirectly irreducibles in V has the amalgamation property, so does V .

## Theorem (F.-Galatos 2022)

The variety of semilinear idempotent residuated lattices has the congruence extension property.

## Centrality

- Let $x^{*}=x \backslash e \vee e / x$. It follows from (F.-Galatos 2022) that if $\mathbf{A}$ is a idempotent residuated chain and $x \in A$, then $x$ fails to commute with at most one element and that element is $x^{*}$.
- Thus $x x^{*}=x^{*} x \Rightarrow x=e$ expresses that the only central element in an idempotent residuated chain $\mathbf{A}$ is $e$.
- We can show that if $w_{S}$ is minimal, then each member of $\mathrm{V}_{S}$ satisfies this quasiequation.
- We will call these exchange-free, and use the same terminology for the corresponding logics.


## Main theorem

We have proven:

## Theorem:

There are continuum-many axiomatic extensions of FL + contraction + mingle + semilinearity with the DIP. Each of these axiomatic extensions is exchange-free and has no non-trivial extensions.

## Part III:

## Further results

## Extensions without the DIP

Leveraging some known failures of amalgamation in varieties of semilinear idempotent residuated lattices, we can also obtain the following:

## Theorem:

There are continuum-many axiomatic extensions of FL + contraction + mingle + semilinearity refuting the exchange rule, but without the DIP.

## Returning to exchange

- If we add exchange back into the picture, the constructions available in the non-commutative case can't be simulated.
- Structural results on commutative idempotent residuated chains, plus application of one-sided amalgamation, gives:


## Theorem:

There are exactly 60 axiomatic extensions of falsum-free FL + exchange + contraction + mingle + semilinearity with the DIP.

- The proof of this amounts to a technical counting argument, not so different from Maksimova's result on intuitionistic logic (hinges on forbidden configurations).


## Adding falsum

- The picture doesn't change that much if we return the falsity constant $f$ to the signature.
- There are still finitely many extensions with DIP in the case with exchange + contraction + mingle + semilinearity.
- But there are many more, and counting them is rather tedious.
- Main idea is that the placement of $f$ in linearly ordered models determines how to decompose these models as a nested sum.
- Combined with the results of (F.-Santschi 2023), this work resolves most of the questions about the number of extensions with DIP for FL + basic structural rules.
- Most interesting open questions involve the weakening rule.
- The tools to get these results are frustratingly diverse, and also require a lot of technology that didn't exist just a few years ago.
- Could pose similar questions about Craig interpolation, uniform interpolation, and so forth.
- This would require new basic tools from, e.g., universal algebra and order-algebraizable logics.


## Thank you!

## Thank you!

